

Decentralized Sequential Hypothesis Testing & Change Detection

Giorgos Fellouris, Columbia University, NY, USA
George V. Moustakides, University of Patras, Greece



Outline

- ★ Sequential hypothesis testing and SPRT
- ★ Sequential change detection and CUSUM
- ★ Decentralized detection and corresponding models
- ★ Centralized schemes (points of reference)
- ★ Decentralized detection using asynchronous random sampling
- ★ Simulation comparisons



Sequential hypothesis testing and SPRT

Conventional binary hypothesis testing (fixed sample size): Collection of observations ξ_1, \dots, ξ_K

$$\mathbb{H}_0: \xi_1, \dots, \xi_K \sim f_0(\xi_1, \dots, \xi_K);$$

$$\mathbb{H}_1: \xi_1, \dots, \xi_K \sim f_1(\xi_1, \dots, \xi_K);$$

Decision rule $D(\xi_1, \dots, \xi_K) \in \{0, 1\}$

$\mathbb{P}(D=1 \mid \mathbb{H}_1)$ (Correct decision)

$\mathbb{P}(D=1 \mid \mathbb{H}_0)$ (Type I error)

$\mathbb{P}(D=0 \mid \mathbb{H}_1)$ (Type II error)

$\mathbb{P}(D=0 \mid \mathbb{H}_0)$ (Correct decision)



Bayes and Neyman-Pearson formulation

Likelihood ratio test:

$$\frac{f_1(\xi_1, \dots, \xi_K)}{f_0(\xi_1, \dots, \xi_K)} \stackrel{\mathbb{H}_1}{\stackrel{\mathbb{H}_0}{\gtrless}} \gamma$$

For i.i.d.:

$$u_K = \sum_{n=1}^K \log \left(\frac{f_1(\xi_n)}{f_0(\xi_n)} \right) \stackrel{\mathbb{H}_1}{\stackrel{\mathbb{H}_0}{\gtrless}} \log(\gamma) = \gamma'$$

WAIT until K samples become available, **THEN** decide

Sequential binary hypothesis testing

Observations $\xi_1, \dots, \xi_n, \dots$ are supplied **sequentially**.

$$\mathbb{H}_0: \xi_1, \dots, \xi_n, \dots \sim f_0(\xi_n)$$

$$\mathbb{H}_1: \xi_1, \dots, \xi_n, \dots \sim f_1(\xi_n)$$

Time Observations

1

ξ_1

Can ξ_1 make a
reliable decision?

2

ξ_1, ξ_2

$\mathcal{N}(\xi_1, \dots, \xi_n) = \{\text{stop}, \text{continue}\}$

...

...

\mathcal{N}

$\xi_1, \dots, \xi_{\mathcal{N}}$

Time \mathcal{N} is
RANDOM

Decision Rule

We **stop** receiving
observations

$$D(\xi_1, \dots, \xi_{\mathcal{N}}) \in \{0, 1\}$$

WHY sequential?

For the same level of confidence with a sequential test we need, in the average, **(significantly) less samples** than a fixed sample size test, to reach a decision.

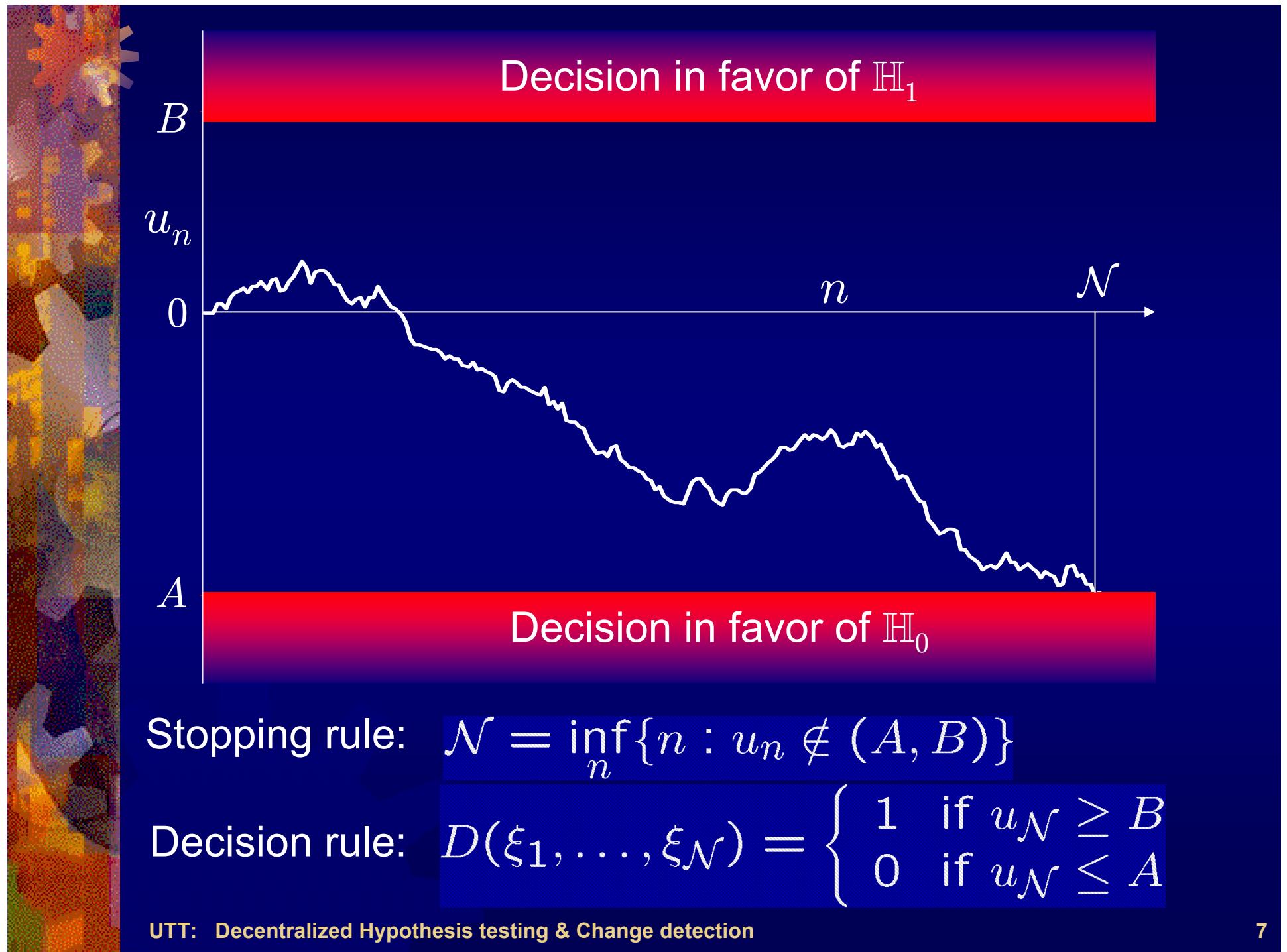
The Sequential Probability Ratio Test (SPRT) (Wald 1947)

Changes with
time

$$u_n = \sum_{k=1}^n \log \left(\frac{f_1(\xi_k)}{f_0(\xi_k)} \right)$$

$$u_n = u_{n-1} + \log \left(\frac{f_1(\xi_n)}{f_0(\xi_n)} \right)$$

We define **two** thresholds $A < 0 < B$



Remarkable optimality property of SPRT

$$\min_{\mathcal{N}, D} \mathbb{E}[\mathcal{N} | \mathbb{H}_0]$$

$$\mathbb{P}[D_{\mathcal{N}} = 1 | \mathbb{H}_0] \leq \alpha; \quad \mathbb{P}[D_{\mathcal{N}} = 0 | \mathbb{H}_1] \leq \beta$$

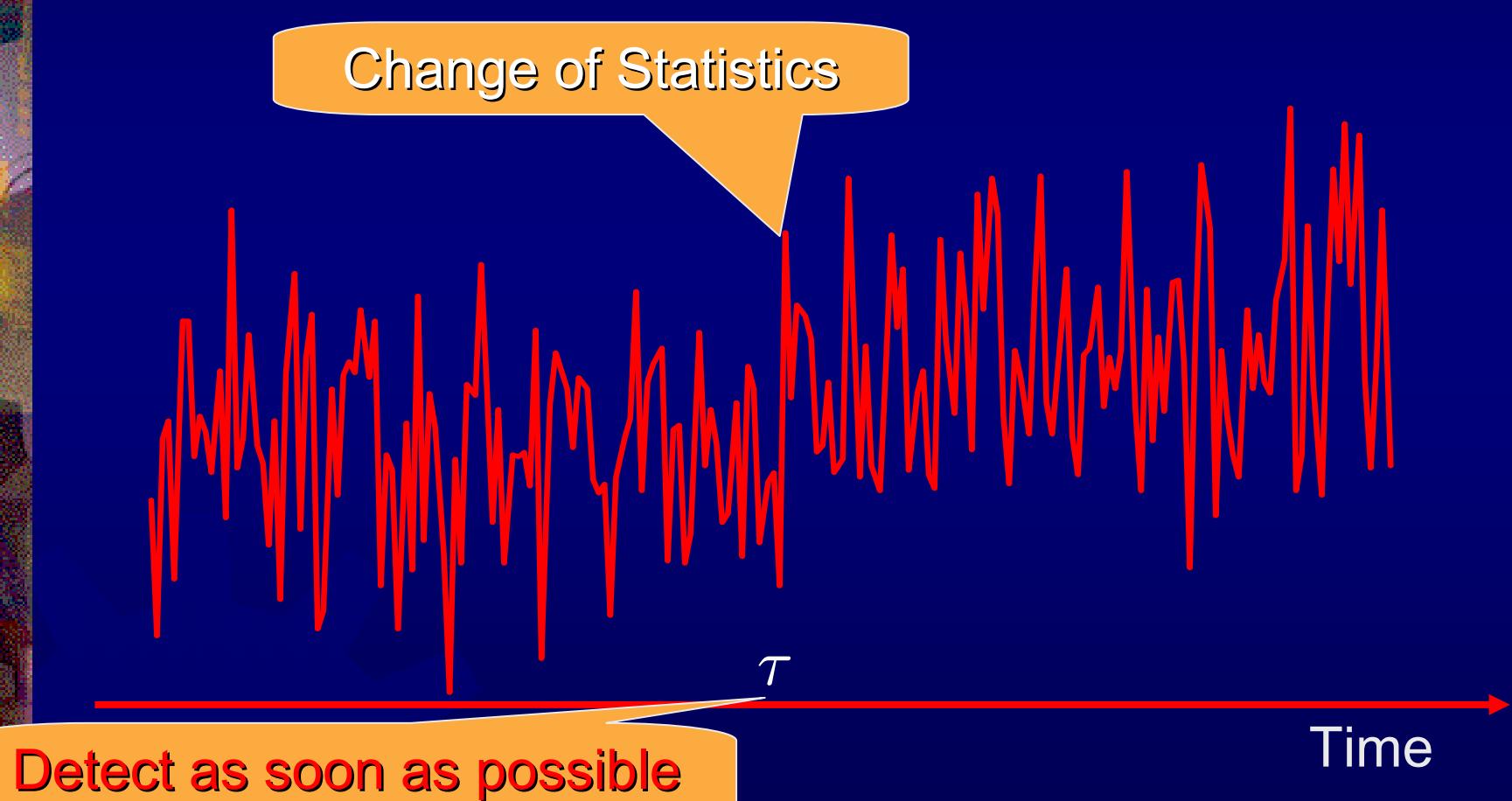
$$\min_{\mathcal{N}, D} \mathbb{E}[\mathcal{N} | \mathbb{H}_1]$$

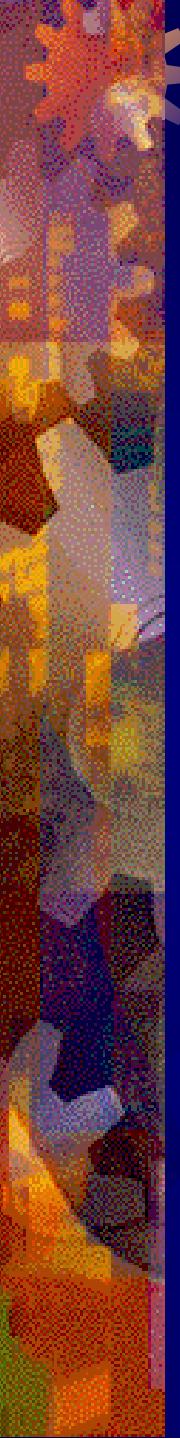
SPRT solves **BOTH** problems **simultaneously**

- * Proved by Wald and Wolfowitz in 1948.

The Sequential change detection problem

Also known as the **Disorder problem** or the **Change-Point problem** or the **Quickest Detection problem**.





Applications

Monitoring of quality of manufacturing process (1930's)

Biomedical Engineering

Electronic Communications

Econometrics

Seismology

Speech & Image Processing

Vibration monitoring

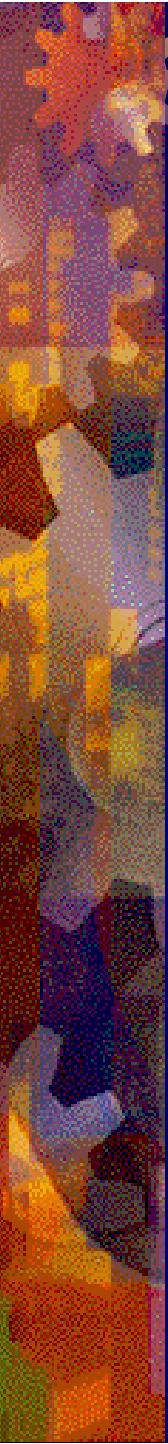
Security monitoring (fraud detection)

Spectrum monitoring

Scene monitoring

Network monitoring and diagnostics (router failures, intruder detection)

Databases



Mathematical setup

We are observing sequentially a process $\{\xi_n\}$ with the following statistics:

$$\begin{aligned}\xi_n &\sim f_0 \quad \text{for } 0 < n \leq \tau \\ &\sim f_1 \quad \text{for } \tau < n\end{aligned}$$

Goal: Detect the change time τ “as soon as possible”

- ★ Change time τ : unknown
- ★ Densities f_0, f_1 : known
- ★ At every time instant n we perform a test and decide whether there was a change or not. In the former case we stop in the latter we continue sampling.
- ★ The test at time n must be based on the **available information up to time n** (and not on any future information).

Cumulative Sum (CUSUM) test

We recall the running log-likelihood:

$$u_n = \sum_{k=1}^n \log \left(\frac{f_1(\xi_k)}{f_0(\xi_k)} \right)$$

The running minimum: $m_n = \inf_{0 \leq s \leq n} u_s$.

Define the CUSUM process y_n :

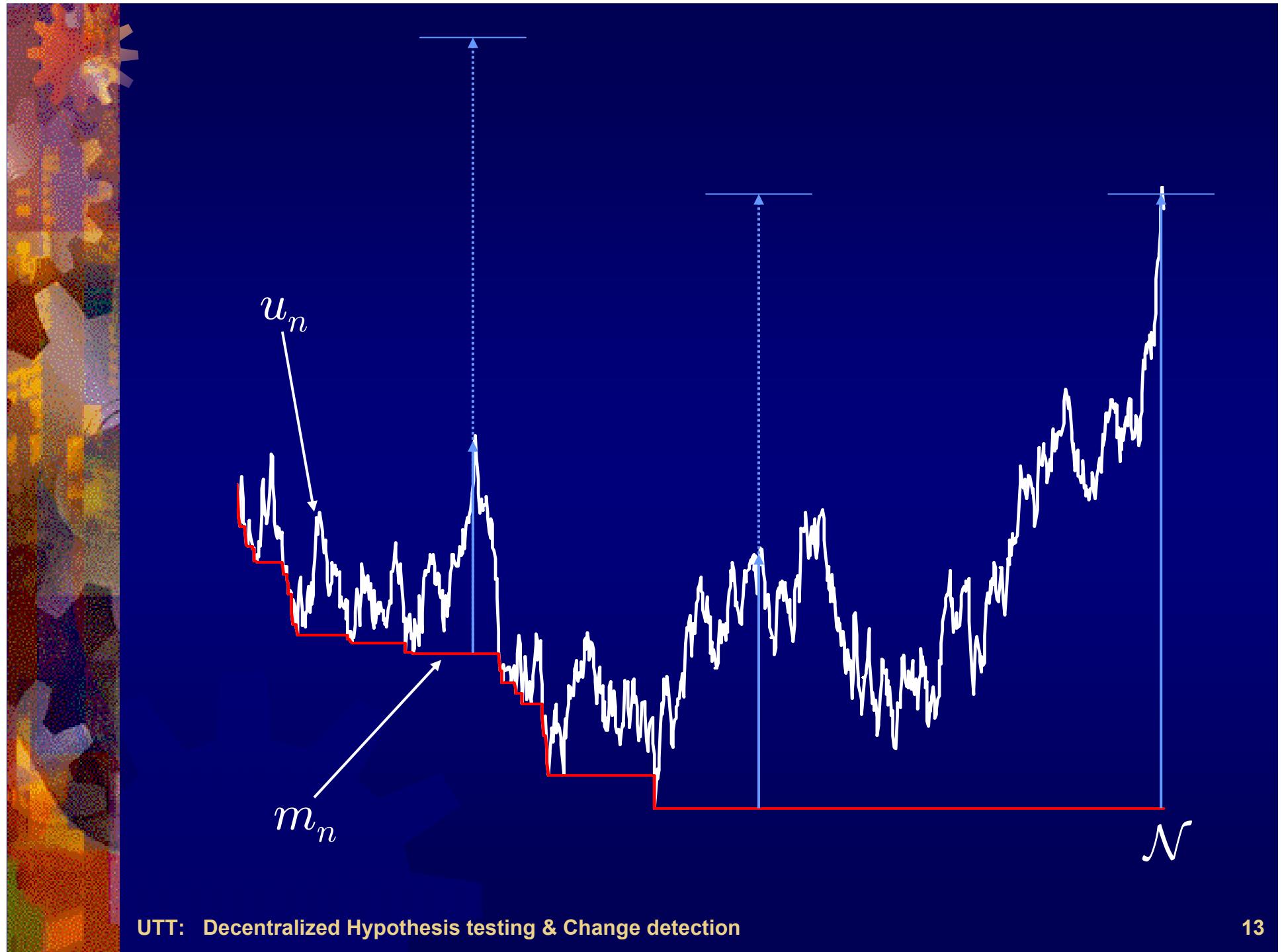
$$y_n = u_n - m_n$$

The CUSUM stopping rule:

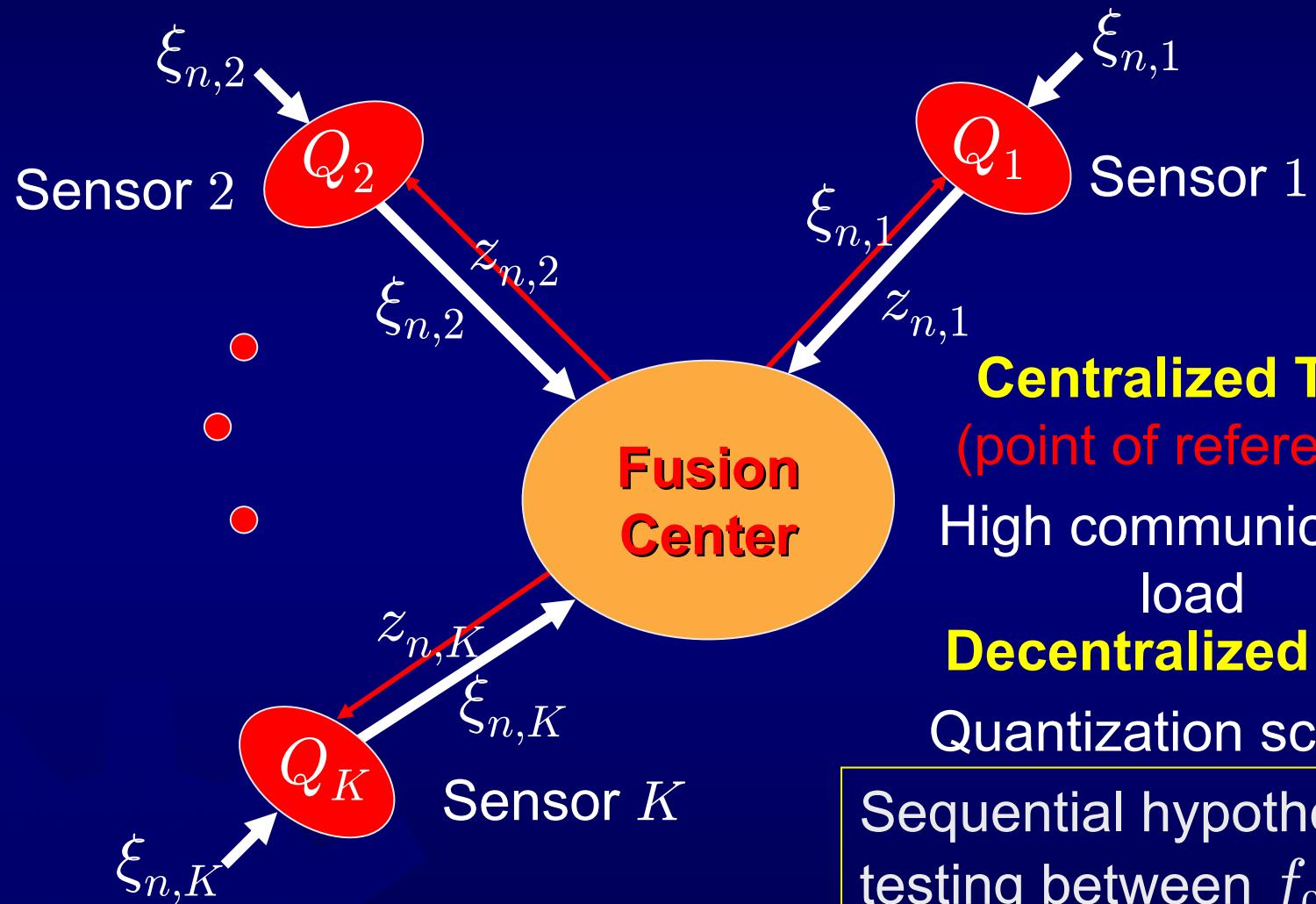
$$\mathcal{N} = \inf_n \{ n: y_n \geq \nu \}$$

We have a convenient recursion:

$$y_n = \max \left\{ y_{n-1} + \log \left(\frac{f_1(\xi_n)}{f_0(\xi_n)} \right), 0 \right\}$$



Decentralized detection and corresponding models



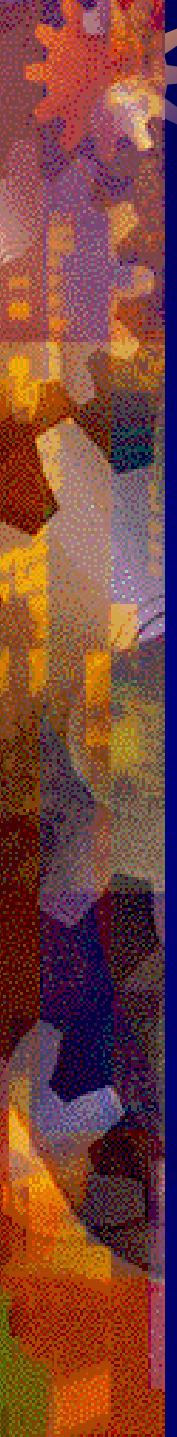
Centralized Test
(point of reference)

High communication
load

Decentralized Test

Quantization scheme

Sequential hypothesis
testing between $f_{0,i}$ and
 $f_{1,i}$.

- 
- ★ No Local Memory:

$$z_{n,i} = Q_i(\xi_{n,i})$$

- ★ Full Local Memory:

$$z_{n,i} = Q_i(\xi_{n,i}, \xi_{n-1,i}, \dots, \xi_{1,i}) \quad \text{Mei (2008)}$$

- ★ Feedback with No Local Memory:

$$z_{n,i} = Q_i(\xi_{n,i}, [z_{n-1,1}, \dots, z_{n-1,K}])$$

- ★ Feedback with Partial Local Memory:

$$z_{n,i} = Q_i(\xi_{n,i}, [z_{n-1,1}, \dots, z_{n-1,K}], \dots, [z_{1,1}, \dots, z_{1,K}])$$

- ★ Feedback with Full Local Memory:

$$z_{n,i} = Q_i(\xi_{n,i}, \dots, \xi_{1,i}, [z_{n-1,1}, \dots, z_{n-1,K}], \dots, [z_{1,1}, \dots, z_{1,K}])$$

Centralized tests

We recall that in this case the sensors send the observations $\xi_{n,i}$ to the Fusion center.

At the Fusion center we form the running log-likelihood ratio

$$\begin{aligned} u_n &= u_{n,1} + u_{n,2} + \cdots + u_{n,K} \\ u_{n,i} &= u_{n-1,i} + \log \left(\frac{f_{1,i}(\xi_{n,i})}{f_{0,i}(\xi_{n,i})} \right) \end{aligned}$$

and apply an SPRT:

Stopping rule: $\mathcal{N}_c = \inf_n \{n : u_n \notin (A_c, B_c)\}$

Decision rule: $D_c(u_1, \dots, u_{\mathcal{N}_c}) = \begin{cases} 1 & \text{if } u_{\mathcal{N}_c} \geq B_c \\ 0 & \text{if } u_{\mathcal{N}_c} \leq A_c \end{cases}$

Remark 1:

In ALL previous detection structures it is assumed the existence of a **GLOBAL CLOCK**.

Synchronization of distant sensors with the fusion center is practically difficult (especially in sensor networks).

Remark 2:

In most practical applications the observation samples $\xi_{n,i}$ come from **canonical sampling of a continuous time process** $\xi_{t,i}$ where

$$\xi_{n,i} = \xi_{nT,i}$$

i.e. we sample $\xi_{t,i}$ at the time instances $t_n = nT$.



An even better centralized scheme

The fusion center instead of receiving the samples $\xi_{n,i}$ it can receive the **CONTINUOUS TIME PROCESSES** $\xi_{t,i}$ to form an SPRT.

$$u_t = u_{t,1} + u_{t,2} + \cdots + u_{t,K}$$

$u_{t,i}$: time evolution of $u_{t,i}$ (sde)

Stopping rule: $\mathcal{T} = \inf_t \{t : u_t \notin (A, B)\}$

Decision rule: $D(u_t, 0 \leq t \leq \mathcal{T}) = \begin{cases} 1 & \text{if } u_{\mathcal{T}} \geq B \\ 0 & \text{if } u_{\mathcal{T}} \leq A \end{cases}$

The continuous time SPRT is **better** than the discrete time SPRT due to infinite time resolution. **It constitutes the ultimate point of reference!**



Asynchronous random sampling

Let t_n^i be increasing sequence of sampling times **NOT necessarily canonical**. At these times we sample the local log-likelihood $u_{t,i}$ in the form $u_{t_n^i, i}$.

Instead of

$$u_t = u_{t,1} + u_{t,2} + \cdots + u_{t,K}$$

we propose the use of the following expression:

$$v_t = u_{t_n^1, 1} + u_{t_n^2, 2} + \cdots + u_{t_n^K, K}$$

Canonical sampling corresponds to: $t_n^i = nT$

Stopping rule: $\mathcal{T} = \inf_t \{t : v_t \notin (A, B)\}$

Decision rule: $D(v_1, \dots, v_{\mathcal{T}}) = \begin{cases} 1 & \text{if } v_{\mathcal{T}} \geq B \\ 0 & \text{if } v_{\mathcal{T}} \leq A \end{cases}$



How do we transmit the local log-likelihoods $u_{t_n^i, i}$ from the sensors to the Fusion center ?

We observe

$$u_{t_n^i, i} = [u_{t_n^i, i} - u_{t_{n-1}^i, i}] + \cdots + [u_{t_1^i, i} - u_{t_0^i, i}]$$

To form the local log-likelihood $u_{t_n^i, i}$ at the fusion center, Sensor i needs to transmit the differences

$$[u_{t_n^i, i} - u_{t_{n-1}^i, i}]$$

We select t_n^i so that the difference $[u_{t_n^i, i} - u_{t_{n-1}^i, i}]$ takes the value A_i or B_i which are specified before hand.

What is the sampling strategy at Sensor i ?

$$t_n^i = \inf_{t > t_{n-1}^i} \left\{ t : [u_{t, i} - u_{t_{n-1}^i, i}] \text{ hits } A_i \text{ or } B_i \right\}$$

- Repeated SPRTs at each sensor.

Communication with the Fusion center

- If the difference is equal to B_i send 1 if it is equal to A_i send 0.
- The Fusion center knows that 1 corresponds to B_i and 0 to A_i and can therefore update the partial log-likelihood then compute v_t and perform the test by comparing v_t to A and B .

- Every time new information arrives at the Fusion center (even from one sensor) the Fusion center updates

$$v_t = u_{t_n^1,1} + u_{t_n^2,2} + \cdots + u_{t_n^K,K}$$

and performs the test.

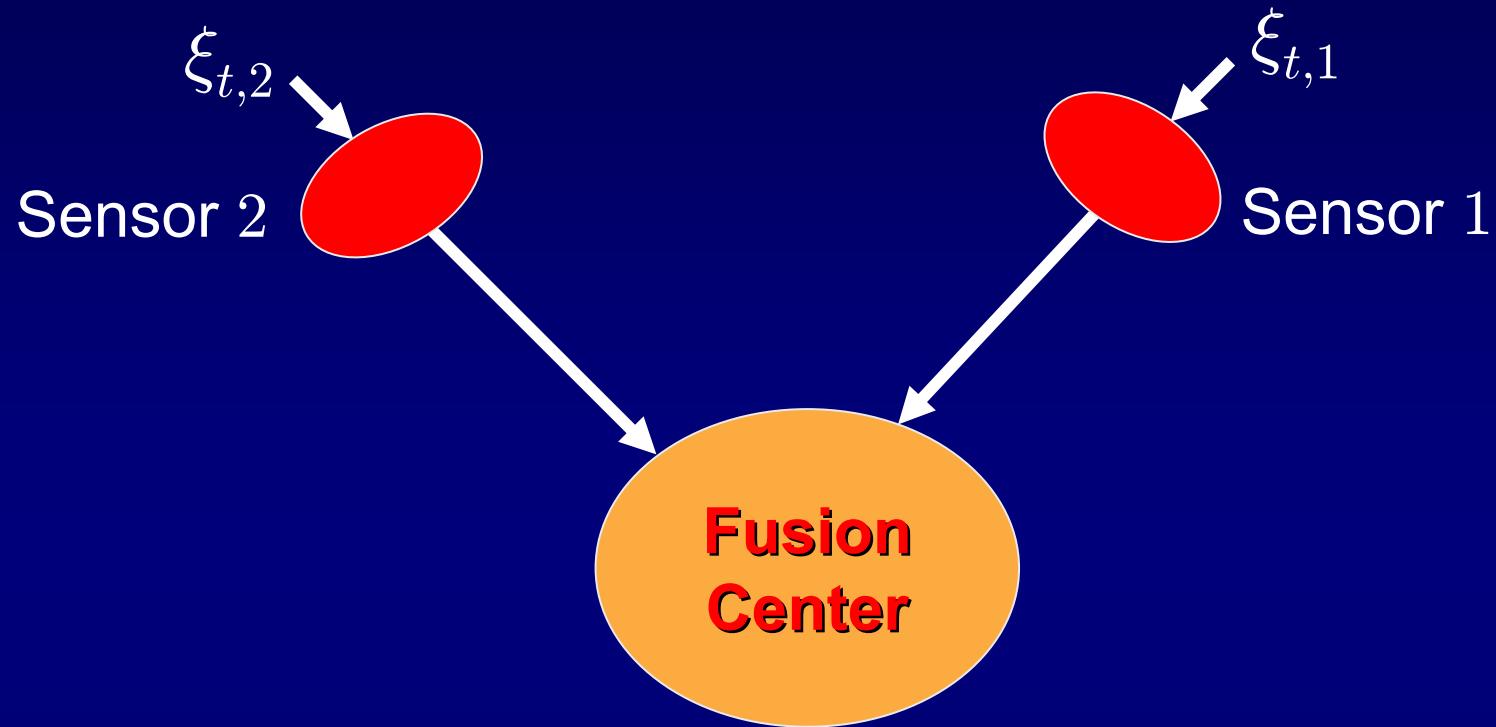
- Communication is Asynchronous and Random!!!

How do we select the local thresholds A_i, B_i ?

We can specify a communication rate between sensors and Fusion center.

If the sensors must communicate, **in the average**, every T time units, then this condition specifies completely the thresholds. We must select the thresholds so that the “average detection delay” of the local SPRTs is equal to T .

Simulations



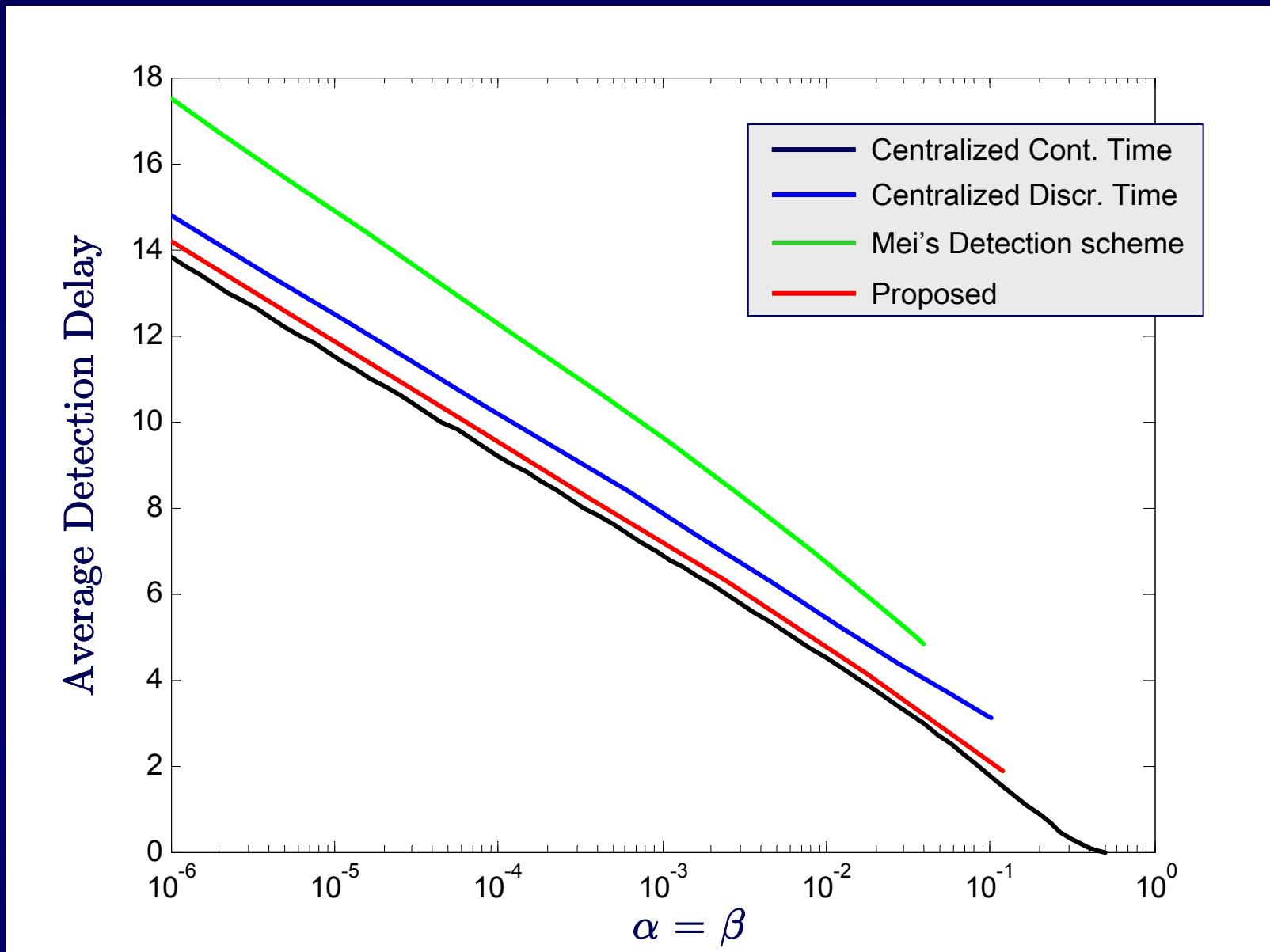
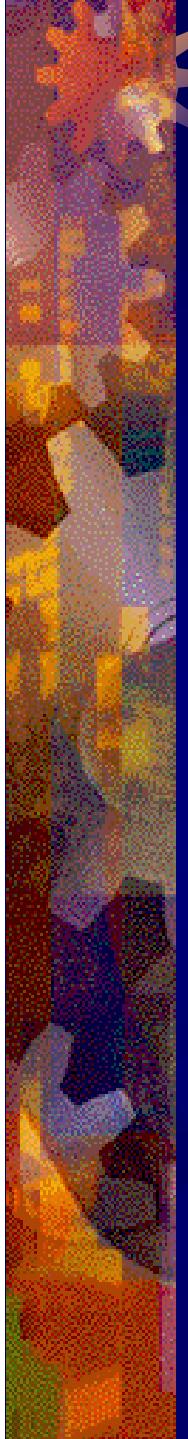
$\xi_{t,i}$: both standard Brownian Motions

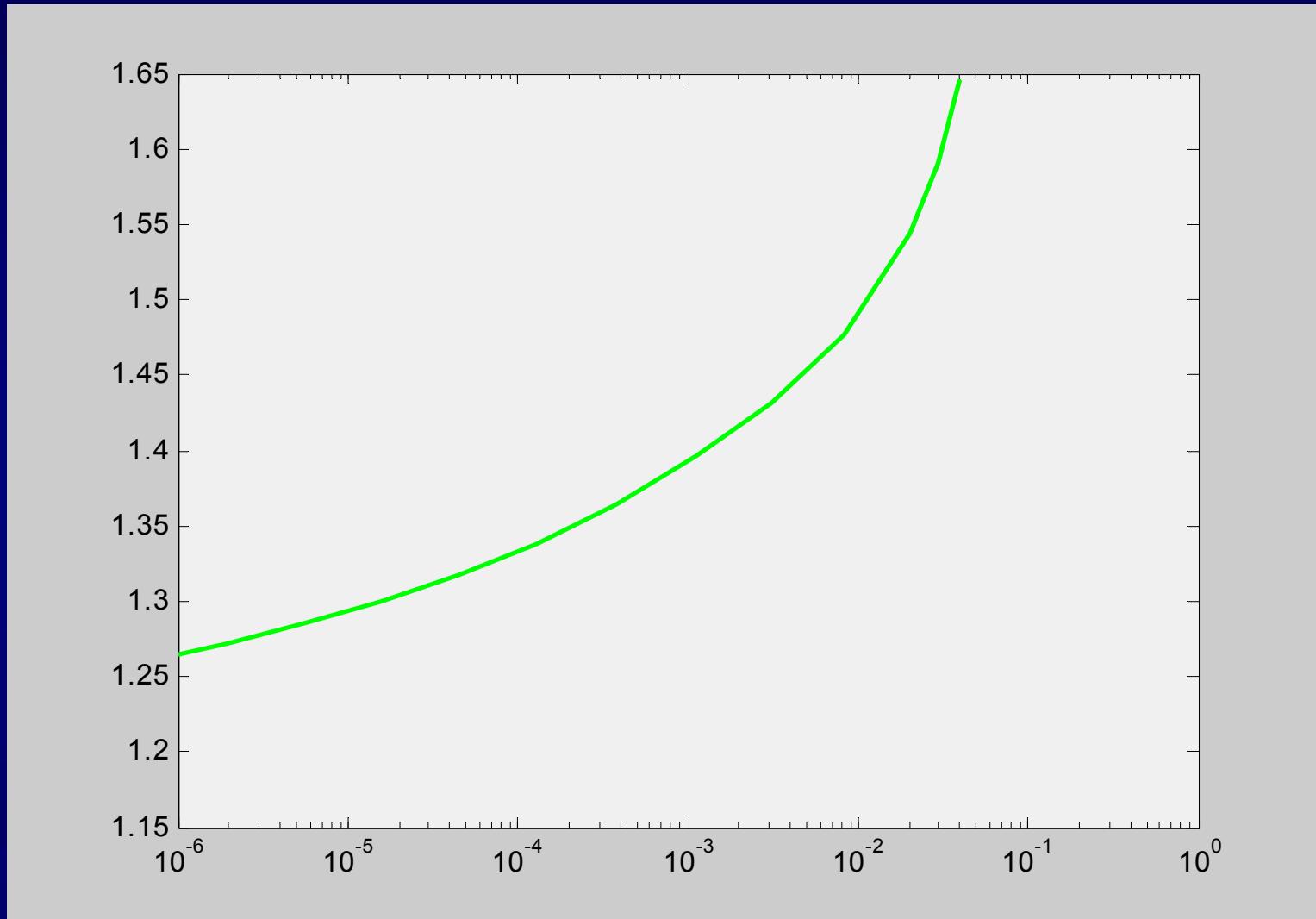
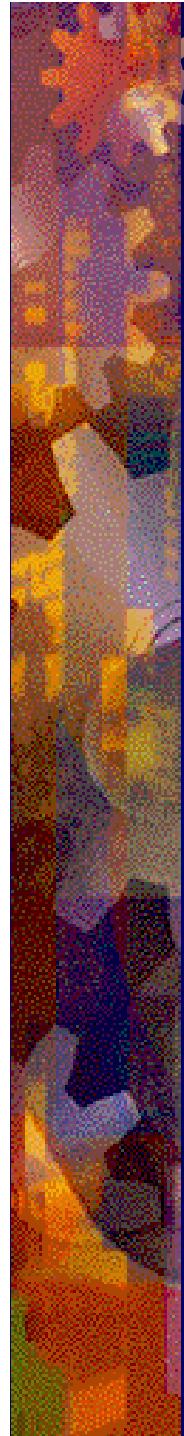
\mathbb{H}_0 : drift = 0; \mathbb{H}_1 : drift = 1.

Average communication period $T=3$.

Canonical sampling generates:

\mathbb{H}_0 : Gaussian(0,3); \mathbb{H}_1 : Gaussian(3,3)







FIN

**Merci de votre
attention**